

# IS-ZC444: ARTIFICIAL INTELLIGENCE

## Lecture-12: Inference with Logical Agents + First Order Logic



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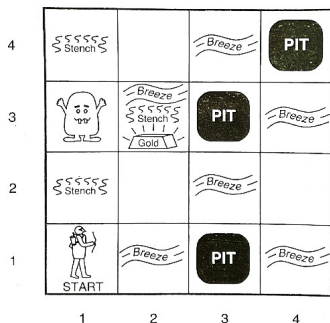
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October 24, 2018

(WILP @ BITS-Pilani Jul-Nov 2018)

# Recall Wumpus World

- **Performance** gold +100, death -100, step -1, arrow -10
- **Environment** smell around wumpus, breeze around pit
- **Actuator** turn left/right, forward, grab, release, shoot
- **Sensor** breeze, glitter, smell, bump, scream



Single Agent, Deterministic, Static, Discrete, !Observable & !Episodic

- $P_{x,y}$  if there is a pit in  $[x, y]$
- $B_{x,y}$  if breeze is in  $[x, y]$
- $W_{x,y}$  if wumpus is in  $[x, y]$
- $S_{x,y}$  if stench is in  $[x, y]$

We know  $R_1: \neg P_{1,1}$ ,  $R_2: B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$ ,  
 $R_3: B_{2,1} \Leftrightarrow (P_{1,1} \vee P_{2,2} \vee P_{3,1})$ ,  $R_4: \neg B_{1,1}$ ,  $R_5: B_{2,1}$

## Model Checking for Inference

- Seven symbols  $P_{1,1}, B_{1,1}, P_{1,2}, P_{2,1}, B_{2,1}, P_{2,2}, P_{3,1}$  have  $2^7 = 128$  models. In three of these knowledge base is true.

$B_{1,1}$	$B_{2,1}$	$P_{1,1}$	$P_{1,2}$	$P_{2,1}$	$P_{2,2}$	$P_{3,1}$	$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$KB$
false	false	false	false	false	false	false	true	true	true	true	false	false
false	false	false	false	false	false	true	true	true	false	true	false	false
⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮
false	true	false	false	false	false	false	true	true	false	true	true	false
false	true	false	false	false	false	true	true	true	true	true	true	<u>true</u>
false	true	false	false	false	true	false	true	true	true	true	true	<u>true</u>
false	true	false	false	false	true	true	true	true	true	true	true	<u>true</u>
false	true	false	false	true	false	false	true	false	false	true	true	false
⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮
true	true	true	true	true	true	true	false	true	true	false	true	false

In all those three  $\neg P_{1,2}$  is true, hence there is no pit in  $[1,2]$ .

On the other hand  $P_{2,2}$  is true on two and false in one so it is not confirmed whether there is pit in  $[2,2]$  or not.

# Validity and Satisfiability

- **Validity:** sentence is true in all models (tautologies)

$$A \vee \neg A$$
$$A \vee B \rightarrow A \vee B$$

- **Satisfiability:** sentence is true in some models

$$A \vee \neg B$$
$$A \rightarrow B$$

Determine whether following sentence is valid or satisfiable

$$((A \wedge B) \rightarrow C) \leftrightarrow (A \rightarrow (B \rightarrow C))$$

# Forward Chaining

Determines if a single proposition symbol  $q$  is entailed by the knowledge? (data driven reasoning)

- It begins from known facts and adds conclusions of the implication whose all the premises are known
- for  $L_{1,1} \wedge breeze \rightarrow B_{1,1}$  if we know  $L_{1,1}$  and *breeze* then  $B_{1,1}$  is added in knowledge base <sup>1</sup>

$$P \Rightarrow Q$$

$$L \wedge M \Rightarrow P$$

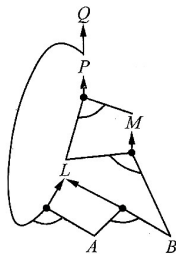
$$B \wedge L \Rightarrow M$$

$$A \wedge P \Rightarrow L$$

$$A \wedge B \Rightarrow L$$

$A$

$B$



- Applies Modus Ponens

$$\frac{\phi \quad \phi \rightarrow \psi}{\psi}$$

<sup>1</sup> $L_{1,1}$ : location is [1,1]

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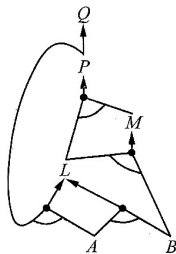
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- Applies Modus Ponens

$$\frac{\phi \quad \phi \rightarrow \psi}{\psi}$$

- An and-or tree gets constructed

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# Backward Chaining

- Works backward from query
- If query  $Q$  is known to be true, then no work is needed.
- Otherwise, find those implications whose conclusion is  $Q$
- If all the premises of one of those implications can be proven true (by backward chaining) then  $Q$  is true

$$P \Rightarrow Q$$

$$L \wedge M \Rightarrow P$$

$$B \wedge L \Rightarrow M$$

$$A \wedge P \Rightarrow L$$

$$A \wedge B \Rightarrow L$$

$A$

$B$

- test( $Q$ ) is it true ?
- test( $P$ ) is it true ?
- test( $L \wedge M$ ) ?
- ((test( $A \wedge B$ ) or test( $A \wedge P$ )) and test( $B \wedge L$ ) ? we know  $A$  and  $B$  so we have  $L$  this gives  $M$
- Therefore  $P$  and hence  $Q$

## First Order Logic (Predicate Logic)

- We have **constants**, **variables**, **predicates** and **functions**
- Here  $P(x)$  could mean  $\forall x$  we have  $P(x)$  or  $\exists x$  such that  $P(x)$
- Variable  $x$  has a domain from where it gets values
- $\forall x, \exists y P(x, y)$  is not always same as  $\exists y, \forall x P(x, y)$
- When we say  $\exists$  a predicate then it is higher order logic

### Examples

- 1 Not every customer have purchased milk and bread



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## Examples

- 1 Not every customer have purchased milk and bread

$$\exists c \text{ Cust}(c) \wedge [\neg \text{shop}(\text{milk}, c) \vee \neg \text{shop}(\text{bread}, c)]$$

- 2 Only one customer have purchased guitar

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$$\exists x [\text{Cust}(x) \wedge \text{shop}(G, x) \wedge \forall y [\neg(x = y) \wedge \text{Cust}(y) \Rightarrow \neg \text{shop}(G, y)]]$$

- 3 Only one customer have purchased guitar and pen

- 4 Highest purchase in forenoon is more than afternoon.

# Inference in First Order Logic

- **Universal Elimination**  $\forall x \text{ Feels}(x, \text{king})$  could be  $\text{Feels}(\text{Raju}, \text{king})$  substitution  $\{x/\text{Raju}\}$  is done using some ground term.
- **Existential Elimination**  $\exists x \text{ Feels}(x, \text{king})$  could be  $\text{Feels}(\text{man}, \text{king})$  if *man* does not appear in knowledge base <sup>2</sup>
- **Existential Introduction** If  $\text{Feels}(\text{Raju}, \text{king})$  then we can say  $\exists x \text{ Feels}(x, \text{king})$

- 1 It is crime for Magadh to sell formula to a hostile country
- 2 Country Bhind, an enemy of Magadh have purchased some formula from Dara
- 3 Dara is from Magadh
- 4 **Question:** Is Dara a criminal?

<sup>2</sup>*man* is a name of person who feels like king

# Prolog

- A logic programming language <sup>3</sup>
- Compile as ['a.pl'].
- If :- and , or ; not **not**
- write('hello'), nl

```
warm_blood(penguin).  
warm_blood(human).  
produce_milk(penguin).  
produce_milk(human).  
have_feather(penguin).  
have_hair(human).  
mammal(X) :-  
warm_blood(X),  
produce_milk(X),
```

```
have_hair(X).
```

- is\_even(X) :-  
Y is X//2, X =:= 2\*Y.
- write('what is your name/ '), read(X), write('Hi '),write(X).

?- mammal(penguin)

no

?- mammal(X).

X = human.

Many more things are possible

<sup>3</sup><http://www.swi-prolog.org/>

# Thank You!

**Thank you very much for your attention!**

**Queries ?**

(Reference<sup>4</sup>)

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<sup>4</sup>1) Book - *AIMA*, ch-07/08, Russell and Norvig. 2) Book - *Logic in CS*, ch-02, Mitchel Huth and Mark Ryan.